- 1. For a dynamical system (T, X), a point $x \in X$ is called *eventually periodic* if there exists m > m' so that $T^m x = T^{m'} x$.
 - (a) Let (T, [0, 1)) be the doubling map and let (σ, Ω) be the full two shift. For each system, give an example of (i) a point that is eventually periodic, and (ii) a point that is eventually periodic but *not* periodic.

Recall, the doubling map is defined by,

$$T: [0,1) \to [0,1); x \mapsto 2x \mod 1.$$

A periodic point can be given by $\frac{1}{5}$ and an eventually periodic point that is not periodic is given by $\frac{11}{24}$:

$$\frac{11}{24}\mapsto \frac{11}{12}\mapsto \frac{5}{6}\mapsto \frac{2}{3}\mapsto \frac{1}{3}\mapsto \frac{2}{3}\mapsto \frac{1}{3}\mapsto \cdots$$

Recall, the full two shift, (σ, Ω) , is defined as, $\Omega = \{0, 1\}^{\mathbb{N}}$ and $\sigma : \Omega \to \Omega$,

 $\sigma(a_0, a_1, a_2 \ldots) = (a_1, a_2, a_3 \ldots).$

Then, a periodic point can be given by the base 2 expansion of $\frac{1}{5} = 0.0011001100110011..._2$. An eventually periodic point that is not periodic is given by the base 2 expansion of the point $\frac{11}{12} = 0.01110101010101..._2$.

$$\sigma^i(p) = \sigma^{i+2}(p), \forall i > 5$$

That is, after shifting 5 times or more, $\frac{11}{12}$ has period of 2.

(b) Let \mathbb{C} be the coding function for the doubling map with the usual partition ($\mathcal{P}_0 = [0, 1/2)$ and $\mathcal{P}_1 = [1/2, 1)$).

Prove $\mathbb{C}(x)$ is eventually periodic in (σ, Ω) if and only if x is eventually periodic in (T, [0, 1)).

Any
$$x \in [0,1]$$
 has a binary expansion given by $x = \sum_{n=0}^{\infty} \frac{x_n}{2^{n+1}}$ for $x_n \in \{0,1\}$.

For $\mathbb{C}(x)$, we may choose the encoding to be the binary values 0 and 1. Then for the doubling map with the defined partitions, $\mathbb{C}(x)$ will encode values x_n to 0 if $T^n x \in \mathcal{P}_0$ otherwise it will encode x_n to 1. So $\mathbb{C} : [0,1) \mapsto \Omega$

Next, we can define $\pi : \Omega \mapsto [0, 1)$ by

$$\pi(x_0, x_1, x_2, \cdots) = \sum_{n=0}^{\infty} \frac{x_n}{2^{n+1}}$$

Then,

$$T(\pi(x_0, x_1, x_2, \cdots)) = T\left(\sum_{n=0}^{\infty} \frac{x_n}{2^{n+1}}\right)$$
$$= x_0 + \sum_{n=1}^{\infty} \frac{x_n}{2^n} \mod 1$$
$$= \sum_{n=1}^{\infty} \frac{x_n}{2^n}.$$

Since $x_0 = 0$

Alternatively,

$$\pi(\sigma(x_0, x_1, x_2, \cdots)) = \pi(x_1, x_2, \cdots)$$
$$= \sum_{n=1}^{\infty} \frac{x_n}{2^n}.$$

Thus, $T \circ \pi = \pi \circ \sigma$. We see that π is surjective and is injective on the set [0, 1) after removing binary expansions ending in all 0s or 1s (dyadic rationals). This is because these rationals all have two binary representations so are not unique.

Since it is trivial to see that dyadic rationals will have period of length 1 in both systems, we may remove them and only consider the remaining values in [0, 1) for which there is a bijection. Since both mappings will have the same binary representation, it is clear that $\mathbb{C}(x)$ is eventually periodic in (σ, Ω) if and only if x is eventually periodic in (T, [0, 1))

(c) Prove that $\mathbb{C}(x)$ is eventually periodic if and only if x is a rational number.

As we have already shown that $\mathbb{C}(x)$ is surjective with the base 2 expansion of a number $x \in [0, 1)$, we know that $\mathbb{C}(x)$ will always correspond to a binary representation of a number in [0, 1). We only need to prove that the binary expansion of any rational number is periodic.

We see that if $x = \frac{1}{n}$ for n > 0 we have $2^m \equiv 1 \mod n$. Then for some $a, 1 \leq a < 2^m - 1$ we have,

$$\frac{1}{n} = \frac{a}{2^k - 1} \\ = \frac{a}{2^k} + \frac{a}{2^{2k}} + \frac{a}{2^{3k}} + \cdots$$

This is equal to the number $x \in [0, 1)$ with base 2 expansion of $0.\overline{00...01}$ where there are k - 1 0s. Thus we see that any rational number will have a periodic binary expansion.

(d) Prove that a binary expansion of a number in [0,1) (viewed as an element of (σ, Ω) is eventually periodic if and only if that number is rational. (*Hint:* $\mathbb{C}(x)$ is always a binary expansion, but not all binary expansions come from codings).

The previous proofs have already shown that all binary expansions $C = \{\mathbb{C}(x) | x \in [0, 1)\}$ will have a periodic expansion iff x is rational. We can see that $\mathcal{B} = \Omega - C$ will be the set of possible binary expansions that do not come from the set of all possible codings.

We find that elements of \mathcal{B} must be irrational because they cannot be represented in the following form which previously represented any periodic base-2 expansion,

$$\frac{1}{n} = \frac{a}{2^k - 1} = \frac{a}{2^k} + \frac{a}{2^{2k}} + \frac{a}{2^{3k}} + \cdots$$

(e) Prove that the base n expansion of a number in [0,1) is eventually periodic if an only if that number is rational.

We can use the following identity,

$$\sum_{n=1}^{\infty} p^{-kn} = \frac{p^{-k}}{1 - p^{-k}} = \frac{1}{p^k - 1}$$

This is equal to the number $x \in [0, 1)$ with base p expansion of $0.\overline{00...01}$ where there are k - 1 0s. Thus any base p expansion of a number in [0, 1) is eventually periodic iff that number is rational

- 2. Let (σ, Ω) be the full two-shift. Let $G \subseteq \Omega$ bet the set of sequences without two ones in a row. Let $X \subseteq \Omega$ be the set of sequences without *three* ones in a row.
 - (a) Show that (σ, X) is a subshift.

We know that X is σ -invariant, that is $\sigma(X) = X$. This is because after shifting any sequence without three ones in row, the resulting sequence will also not contain three ones in a row.

We also know that X is a closed set because it contains all its limit points. This is true because infinite applications of sigma can not produce sequences containing three consecutive ones in a row if the original sequence did not contain three consecutive ones.

Thus, (σ, X) is a subshift.

(b) A *two-step* Markov chain is a Markov chain where the transition probabilities between states depend on the current state *and* the previous state.

A two-step Markov chain on a state space S can be thought of as a one-step (regular) Markov chain on the state space $S \times S$. For example, let \mathcal{M} be the Markov chain on states 0 and 1 with equal transition probabilities. Normally, \mathcal{M} has a graph like this:



However, we can also model \mathcal{M} as a Markov chain \mathcal{M}' with state space 00, 01, 10, and 11 and transition graph like this:



Explain how \mathcal{M}' models \mathcal{M} . In particular, explain why the graph for \mathcal{M}' has no edge between 00 and 11.

We know the states of \mathcal{M}' capture the previous and current state of \mathcal{M} . \mathcal{M}' can be understood as being a model of \mathcal{M} if we consider $00 \equiv 0, 11 \equiv 1$ and suppose the additional transitions for the states 01 and 10 have equivalent steady state probabilities for $0 \to 1, 1 \to 0$ and self transitions $0 \to 0, 1 \to 1$ respectively. Moreover, the additional states add extra information about the transition probabilities given the previous state.

There is no edge between 00 and 11 because it is not possible for a state 0 to transition to state 1 without having a previous state of 0.

(c) Let \mathcal{M}_G be a Markov chain whose set of realizations is G.

i. Draw a graph for \mathcal{M}_G and find its associated transition matrix.

$$\mathbf{O} \qquad \mathbf{M}_G = \begin{bmatrix} 0.5 & 1 \\ 0.5 & 0 \end{bmatrix} = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}$$

ii. Use the transition matrix for \mathcal{M}_G to compute the entropy of (σ, G) .

We may use the fact that for subshifts of finite type, topological entropy coincides with the logarithm of the spectral radius of the transition matrix¹.

Let λ^* be the spectral radius, i.e. the largest eigenvalue modulus of the adjacency matrix. Then we can calculate the topological entropy as,

$$\mathcal{H}(G) = \lim_{n \to \infty} \frac{\log(\# \text{ words of length n in } G)}{\log(\# \text{ of words of length n in } \Omega)}$$
$$\equiv \log(|\lambda^*|) = \log\left(\left|\frac{(1+\sqrt{5})}{2}\right|\right) = 0.4812118250...$$

iii. Model \mathcal{M}_G as a two-step Markov chain \mathcal{M}'_G , which can be viewed as a Markov chain with state space $\{00, 01, 10, 11\}$. Draw the graph associated with \mathcal{M}'_G and find its transition matrix.



iv. Using the transition matrix for \mathcal{M}'_G , compute the entropy of (σ, G) .

Let λ^* be the spectral radius, i.e. the largest eigenvalue of the adjacency matrix. Then we can calculate the topological entropy as

$$\mathcal{H}(G) = \lim_{n \to \infty} \frac{\log(\# \text{ words of length n in } G)}{\log(\# \text{ of words of length n in } \Omega)}$$
$$= \log(|\lambda^*|) = \log\left(\left|\frac{(1+\sqrt{5})}{2}\right|\right) = 0.4812118250...$$

(d) Can X be modeled by a one-step Markov chain? Why or why not?

Recall, X is the subset of sequences without 3 consecutive 1's. The one-step Markov chain could not model this as removing a transition or edge from 1 to itself would also remove all sequences with two consecutive 1s.

 $^{^{1}} http://www.scholarpedia.org/article/Topological_entropy\#Topological_entropy_in_some_special_cases$

(e) Find the entropy of (σ, X) .

We see that $X \subset G$ without any transitions or edges from state 11 to itself so it can be modeled with the two-state markov chain and its entropy found,



- 3. Let (σ, Ω) be the full two shift.
 - (a) Prove that (σ, Ω) is expansive.

Recall, (σ, Ω) is expansive if for any $x, y \in \Omega$, if $x \neq y$ then for some n and some $\epsilon > 0$, $d(\sigma^n(x), \sigma^n(y)) > \epsilon$.

Suppose $(a_n), (b_n) \in \Omega$ and $(a_n) \neq (b_n)$, then for some $K \in \mathbb{N}, a_K \neq b_K$. Therefore $d(\sigma^K((a_n)), \sigma^K((b_n))) \geq 1$. So the full two shift is an expansive dynamical system.

(b) Prove that (σ, Ω) is *transitive*.

Recall, the point $x \in \Omega$ is called transitive if $\mathcal{O}(x) = \Omega$, where $\mathcal{O}(x)$ the closure of the orbit. The system (σ, Ω) is transitive if there exists a transitive point in Ω .

We claim that there exists a $y \in \Omega$ such that $\overline{\mathcal{O}(y)} = \Omega$. Then, for any $\alpha \in \Omega$, we can choose an arbitrary $\epsilon > 0$, and choose *n* such that $2^{-n} < \epsilon$. Since α will contain some finite *n* symbols and since *y* contains all lists of finite symbols, the contents of α must be in *y* at some index *j*. Then, by shifting *y* to this index *j* will reduce the distance to any arbitrary value, i.e. $d(\sigma^k(y), \alpha) < \epsilon$ for some *k*. We can see that this is the definition for *y* having a dense orbit, which means that (σ, Ω) is transitive.

We may choose the base 2 expansion of some transcendental irrational to be y, then we know it will be neither periodic nor eventually periodic and will contain all possible sequences of binary values.

(c) Show that (σ, Ω) is *chaotic*.

By our definition, the dynamical system (σ, Ω) is chaotic if it is both expansive and transitive.

(d) Let (T, X) be a dynamical system and let $x, y \in X$. We say x and y are δ -correlated for n steps if the distance between $T^i x$ and $T^i y$ is at most δ for i = 0, ..., n.

Let $x, y \in \Omega$ be points that are at distance 2^{-k} of each other. For how many steps will x and y be δ -correlated when $\delta = 1/4$?

If $d(x, y) = 2^{-k}$ we know $k = \min_k \{j : x_j \neq y_j\}$ is the number of steps x and y will be δ -correlated. We also have that $d(T^i x, T^i y) \leq \frac{1}{4} = 2^{-2}$. Thus, x and y be δ -correlated for 2 steps.

(e) If (T, X) is a chaotic dynamical system, can two points be δ -correlated indefinitely? What implications does this have for measurement error?

For two points to be δ -correlated indefinitely means that the distance between $T^n x$ and $T^n y \leq \delta$ for all $n \in \mathbb{N}$.

If X is finite, we may choose some δ outside the radius of the system (T, X) so that two points will be δ -correlated indefinitely.

However, if (T, X) is infinite and chaotic or we pick a point within the radius then, because the system is expansive we know there exists $\epsilon > 0$, so that if $x, y \in \Omega$ and $d(T^n(x), T^n(y)) \leq \epsilon$ for every $n \in \mathbb{N}$ then, x = y. So two points $x \neq y$ will not be δ -correlated indefinitely in an infinite choatic dynamical system or for some choices of δ inside the bounds of a finite system.

If you can always pick a small enough δ so that nothing is delta correlated then you will have a minimum bound on your measurement error. Conversely, if you pick a large enough δ outside the system, you have a maximum bound on your measurement error. 4. Two dynamical systems (T, X) and (S, Y) are called *conjugate* if there exists a continuous, invertible function $\Phi: Y \to X$ so that $T = \Phi^{-1} \circ S \circ \Phi$. In this case, Φ is called a *conjugacy*.

The systems are called *semi-conjugate* if there exist continuous, onto function $\Phi: Y \to X$ so that $T \circ \Phi = \Phi \circ S$. In this case, Φ is called a *semi-conjugacy*.

- (a) Let (T, X) and (S, Y) be conjugate dynamical systems.
 - i. Show that if there exists a point of period k in Y, there exists a point of period k in X.

Since (T, X) and (S, Y) are conjugate, we have,

$$T = \Phi^{-1} \circ S \circ \Phi$$
$$T^{k} = \Phi^{-1} \circ S^{k} \circ \Phi$$
$$S^{k} = \Phi \circ T \circ \Phi^{-1}$$

Then for $y \in Y$,

$$y = S^{k}y$$
$$\implies \Phi^{-1} \circ \Phi y = \Phi^{-1} \circ S^{k} \circ \Phi y$$
$$y = T^{k}y.$$

Since Φ is bijective we know there will be a unique $x \in X$,

$$y = \Phi^{-1}(x)$$
$$\implies x = T^k x$$

Therefore if a point in Y has period k, a point in X will also have period k.

ii. Show that if (S, Y) is transitive, then (T, X) is transitive.

If (S, Y) is transitive, then for some $y \in Y$ we have,

$$\overline{\mathcal{O}(y)} = Y$$
$$= \bigcap_{n \in \mathbb{N}} \{ S^k(y) : k > n \}.$$

Since Φ forms a bijection between Y and X, for each point in Y from the closure of the orbit of y there will be a unique point in X

$$\overline{\mathcal{O}(y)} = \overline{\mathcal{O}(\Phi^{-1} \circ \Phi(y))} = \bigcap_{n \in \mathbb{N}} \{\Phi^{-1} \circ S^k \circ \Phi(y) : k > n\}$$
$$= \bigcap_{n \in \mathbb{N}} \{T^k(y) : k > n\}$$
$$= X$$

Thus, if (S, Y) is transitive, then (T, X) is also transitive

iii. Is it true that if (S, Y) is expansive, then (T, X) is necessarily expansive?

If X and Y are compact, then Φ is uniformly continuous so there exists an $\alpha > 0$ such that

$$d(\Phi(y_1)), \Phi(y_2)) \le \alpha \implies d(y_1, y_2) < \epsilon$$
$$d(y_1, y_2) > \epsilon \implies d(\Phi(y_1), \Phi(y_2)) > \alpha$$

If $x_1, x_2 \in X$, $x_1 \neq x_2$, then $\Phi^{-1}(x_1) \in Y$ and $\Phi^{-1}(x_2) \in Y$ so there exists an $n \in \mathbb{Z}$ such that

$$\begin{aligned} &d(S^n \Phi^{-1}(x_1), S^n \Phi^{-1}(x_1)) > d \\ &d(\Phi^{-1}S^n(x_1), \Phi^{-1}S^n(x_1)) > d \\ &d(\Phi^{-1}S^n \Phi(x_1), \Phi^{-1}S^n \Phi(x_1)) > \alpha \\ &d(T^n x_1, T^n x_2) > \alpha \end{aligned}$$

So, (T, X) is also expansive.

However, if X or Y are not compact and are infinite or open, then the implication doesn't necessarily hold and we can find a counter example as follows

$$T : \mathbb{R} \mapsto \mathbb{R}, T(x) = 2x$$

$$S : [0, 1) \mapsto [0, 1), S(x) = \frac{1}{2x}$$

$$\Phi : T \mapsto S, \Phi(x) = \frac{1}{x}$$

$$\Phi^{-1} \circ S \circ \Phi(x) = \Phi \circ S\left(\frac{1}{x}\right) = \Phi\left(\frac{1}{2x}\right) = 2x$$

Here we see two systems are conjugate by Φ , but T is an expansion while S is contraction. So the statement doesn't hold when X,Y are not compact or when Φ is not uniformly continuous

(b) Prove that if (T, X) and (S, Y) are conjugate dynamical systems, that they are also semi-conjugate dynamical systems.

Since Φ is bijective, it will be surjective. We can derive the identity from the definition as follows,

$$T = \Phi^{-1}S\Phi$$
$$\Phi T = \Phi \Phi^{-1}S\Phi$$
$$\Phi T = S\Phi$$
$$T\Phi = \Phi S$$

Using Φ 's continuous invertibility, it can be shown that function composition will be commutative. Then, this result matches the definition of being topologically semi-conjugate.

- (c) Let (T, X) and (S, Y) be semi-conjugate dynamical systems.
 - i. Show that if there exists a point of period k in Y, there exists a point of period $\leq k$ in X.

It can easily be shown through induction that,

$$T^n \circ \Phi = \Phi \circ S^n$$

If y is a periodic point of period n for S we have,

$$y = S^n y$$

By the conjugacy relation,

$$T^n(\Phi(y)) = \Phi(S^n(y)) = \Phi(y)$$

This means that $\Phi(y)$ is a periodic point of period n for (T, X).

$$\Phi y = x_i$$
$$x_i = T^k x_i$$

Because Φ is surjective,

$$dim(S) \le dim(T)$$
$$k \le n$$

Thus, we can conclude that if there exists a point of period n in Y, there exists a point of period $k \leq n$ in X.

ii. Show that if (S, Y) is transitive, then (T, X) is transitive.

If (S, Y) is transitive, we know there $\mathcal{O}_S(y)$ is dense for some $y \in Y$.

For any non-empty open set $U \subset X$, it must be true that $\Phi^{-1}(U)$ is an open set in Y because Φ is surjective and its inverse is continuous. Using the assumed density of the orbit of y, there must be some $k \in \mathbb{N}$ such that,

$$S^k(y) \in \Phi^{-1}(U) \iff \Phi(S^k(y)) \in U.$$

Because Φ defines a conjugcacy,

$$T^{k}\Phi = \Phi S^{k}T^{k}(\Phi(y)) = \Phi(S^{k}(y)) \in U \implies \mathcal{O}_{T}(\Phi(y)) \cap U$$

This means that if the orbit $\mathcal{O}_S(y)$ is dense in Y, then the orbit $\mathcal{O}_T(\Phi(y))$ is dense in X. Moreover, it implies that (T, X) will contain a transitive point, so it must also be transitive.

iii. Is it true that if (S, Y) is expansive, then (T, X) is necessarily expansive?

Since the semi-conjugacy Φ is only surjective, it won't necessarily commute with S^n . Thus we can not show that (S, Y) being expansive implies (T, X) is expansive in general.

- (d) Let (T, [0, 1)) be the doubling map and let (σ, Ω) be the full two shift.
 - i. Show that (T, [0, 1)) and (σ, Ω) are semi-conjugate.

As seen before, we may claim $\pi : \Omega \mapsto [0,1]$ defines a surjective semiconjugative map,

$$\pi(a_0, a_1, a_2, \cdots) = \sum_{n=0}^{\infty} \frac{a_n}{2^{n+1}}$$

We know that π is surjective since each real $x \in [0, 1]$ has a binary expansion. To show that it commutes so that $\pi \circ \sigma = T \circ \pi$. First we find $\pi \circ \sigma$

$$\pi(\sigma(a_i)_{i=1}^{\infty}) = \pi(\sigma(b_i)_{i=1}^{\infty}) = \sum_{i=1}^{\infty} \frac{b_i}{2^i} = \sum_{i=1}^{\infty} \frac{a_{i+1}}{2^i}$$

Given that $b_i = a_{i+1}$. Next we find $T \circ \pi$

$$T(\pi(a_i)_{i=1}^{\infty})) = T\Big(\sum_{i=1}^{\infty} \frac{a_i}{2^i}\Big) = \sum_{i=1}^{\infty} \frac{a_i}{2^i} \mod 1 = \sum_{j=1}^{\infty} \frac{a_{j+1}}{2^j}$$

We see the results are the same, so $\pi \circ \sigma = T \circ \pi$ and π defines a semi-conjugacy between (T, [0, 1)) and (σ, Ω) .

ii. Show that (T, [0, 1)) is chaotic.

We have already shown that (σ, Ω) is chaotic and thus transitive. Since we've shown that (T, [0, 1)) is semi-conjugate to (σ, Ω) with our defined π we know that it is transitive, then we only need to show that (T, [0, 1)) is expansive.

Let $x \neq y$ be distinct points in [0, 1), if $d(x, y) < \frac{1}{4}$ we see that

$$d(T(x), T(y)) > 2d(x, y)$$

Then for all n and for some $\epsilon > 0$

$$d(T^n(x), T^n(y)) = 2^n d(x, y) > \epsilon$$

Thus, (T, [0, 1)) is expansive and transitive so it must be chaotic. Moreover, we found that the expansive constant is $\epsilon = \frac{1}{4}$.

- (e) Let (T, [0, 1)) be the doubling map and let (L, [0, 1]) be the *logistic map* defined by L(x) = rx(1-x) for r = 4.
 - i. Define $f : [0,1) \to [0,1)$ by $f(x) = \sin^2(2\pi x)$. Show that f is a semi-conjugacy between (L, [0,1]) and (T, [0,1)).

We may first note that,

$$f(x) = \sin^2(2\pi x) = \frac{1}{2}(1 - \cos(4\pi x))$$

We want to show that f is surjective and commutes as expected between T and L so we need to verify that $L \circ f = f \circ T$. We first find $L \circ f$,

$$L \circ f = 4f(x)(1 - f(x))$$

= 2(1 - cos(4\pi x)) $\frac{1}{2}(1 + cos(4\pi x))$
= 1 - cos²(4\pi x).

Next, we find $f \circ T$,

$$f \circ T = \frac{1}{2}(1 - \cos(4\pi T(x)))$$
$$= \frac{1}{2}(1 - (2\cos^2(4\pi x) - 1))$$
$$= 1 - \cos^2(4\pi x).$$

Thus, $L \circ f = f \circ T$. We see that f is a composition of $cos(4\pi x)$ and $\frac{1-x}{2}$, both of which are continuous and surjective on [0, 1) so their composition will also be surjective and continuous. Then f gives a semi-conjugacy.

ii. Show that (L, [0, 1]) contains points of every period.

We know the logistic map is semi-conjugate to the doubling map. We also know the doubling map is semi-conjugate to the full two-shift map. Then, we try to define a semi-conjugacy g between (L, [0, 1]) and (σ, Ω) with

$$g(x) = \sin^2(2\pi(x)).$$

However, we see that for points of the form, $x \in [0, 1), x = \frac{1}{q}$ where $2 \nmid q, x$ will have period of q - 1 but g will collapse half of those values along their symmetry of 1/2 of the sine curve. We may claim that for these points, $g(\frac{1}{q})$ has a period of q - 1 in T but will have half that period of $\frac{q-1}{2}$ in L.

We see that the sorted values of the orbit will be $\mathcal{O}(\frac{1}{q}) = \{\frac{1}{q}, \frac{2}{q}, \frac{3}{q}, ..., \frac{q-1}{q}\}$ and will have symmetry along 1/2. There will be an even number of equally spaced points in the period with half in the LHS and the other half in the RHS. If we consider only the first half from [0, 1/2] we see the right endpoint will have distance of 1/2q to 1/2, but the left point will have distance 1/qfrom 0. This means that the points can not be symmetrically distributed along 1/4 and moreover, it means that the points we be guaranteed to have no overlap in g.

We have already seen how periods of any size k can be formed from the full two shift map using an identity for rationals. Form here, it is easy to see how these rationals can be constructed so that they are mapped to a value in (L, [0, 1]) with arbitrary period lengths using the above defined conjugacy.

iii. Show that (L, [0, 1]) is chaotic.

It can be shown that $g(x) = \sin^2(2\pi(x))$ defines a conjugate bijective map between (L, [0, 1]) and (σ, Ω) .

We've shown that (σ, Ω) is chaotic, then since (L, [0, 1]) is conjugate, it must be transitive and expansive as well and is therefore also chaotic.

Programming Problems

For the programming problems, please use the Jupyter notebook available at

https://utoronto.syzygy.ca/jupyter/user-redirect/git-pull?repo=https://github.com/siefkenj/ 2020-MAT-335-webpage&subPath=homework/homework4-exercises.ipynb

Make sure to comment your code and use "Markdown" style cells to explain your answers.

1. The logistic map with parameter r is the function $f:[0,1] \rightarrow [0,1]$ defined by f(x) = rx(1-x)(for $r \in [0,4]$). It was invented as a simple model for population in biology. The idea is that if x is your original population, there will be some birth rate r governing population growth; however, when the population outstretches its resources, its growth will be constrained. The 1-x parameter models this constraint on growth.

Despite being so simple, the logistic map can have amazingly complex behaviour!

- (a) A two-parameter logistic function f has been predefined. Create a function orbit_f which inputs a starting value, a parameter r, and an orbit length n, and returns a list with the first n points along the orbit of f.
- (b) Given a dynamical system (T, X), the *n*-orbit of x is $\mathcal{O}^n(x) = \{x, Tx, \dots, T^{n-1}x\}$.

Using r = 2, plot the 20-orbits of x for at least 10 different x's (plot time vs. value). What do you notice?

- (c) Does the logistic map with r = 2 have a basin of attraction²? Justify your conclusion.
- (d) Repeat part 1b using r = 3. What do you notice? Is there a basin of attraction consisting of a single point?
- (e) In general, a basin of attraction for a dynamical system is a set which all points limit to. Does the logistic map have a basin of attraction when r = 3? If so, describe it.
- (f) Repeat part 1b with r = 4. Is there a basin of attraction? Why or why not?
- 2. We are going to plot the basins of attraction for the logistic map as a function of r. We can approximate a basin of attraction by taking the 1000-orbit of several points, and then taking the set consisting of the last 100 points in each orbit.
 - (a) Create a function approximate_basin which takes in a parameter r and returns a list of points that approximate the basin of attraction of the logistic map with parameter r.

Hint: It may be useful to round your results to 3 or 4 decimal places and then use np.unique to get a list of manageable length.

- (b) Plot the basins of attraction of the logistic map vs. r for at 1000 different values of r between 0 and 4.
- (c) What do you notice about the basins of attraction? Did you expect this?
- (d) In the *proofs* part of this homework set, you proved that the logistic map is chaotic when r = 4. What does that imply about the basin of attraction?
- (e) If a population is well-modeled by the logistic map with a parameter of r = 3.82, what can you say about the population after 1000 days? What if it is modeled by a logistic map with parameter r = 3.83? Is the population more or less predictable?

²Look back in the notes if you've forgotten this term.